

Averaging over local unitary groups in Random Tensor Networks

Sylvain Carrozza



Random tensors and related topics
Institut Henri Poincaré
October 18, 2024

OUTLINE

- ▶ Generalities on entanglement and local unitary invariants
- ▶ Review of standard results on RTNs
- ▶ RTNs with reduced randomness: local Haar-averaging

[Work in progress with Luca Lionni + ...]

Generalities on entanglement and local unitary invariants

SEPARABLE AND ENTANGLED STATES

$$\mathcal{H} = \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \cdots \otimes \mathcal{H}_q$$

Pure states. We say that $|\psi\rangle \in \mathcal{H}$ is

- ▶ *separable* if $\exists\{ |v_i\rangle \}$ such that

$$|\psi\rangle = |v_1\rangle \otimes |v_2\rangle \otimes \cdots \otimes |v_q\rangle$$

- ▶ *entangled* otherwise.

Mixed states. ρ is

- ▶ *separable* if $\exists\{\rho_i^{(k)}\}$ and $\{\alpha_i | \alpha_i \geq 0, \sum_i \alpha_i = 1\}$ such that

$$\rho = \sum_k \alpha_k \rho_1^{(k)} \otimes \rho_2^{(k)} \otimes \cdots \otimes \rho_q^{(k)}$$

- ▶ *entangled* otherwise.

LOCAL UNITARY SYMMETRY

$$\mathcal{H} = \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \cdots \otimes \mathcal{H}_q$$

Local Unitary (LU) transformation. Change of frame in each local subsystem \mathcal{H}_i , represented by

$$U_1 \otimes U_2 \otimes \cdots \otimes U_q \quad \text{with} \quad U_i \in U(D_i) \quad \forall i$$

Entanglement structure. Orbit under action of $U(D_1) \otimes U(D_2) \otimes \cdots \otimes U(D_q)$ i.e. equivalence class of states:

► *pure states:*

$$|\psi\rangle \sim_{\text{LU}} |\psi'\rangle \quad \Leftrightarrow \quad |\psi'\rangle = (U_1 \otimes U_2 \otimes \cdots \otimes U_q) |\psi\rangle$$

Example. {separable states} constitute one such equivalent class.

► *mixed states:*

$$\rho \sim_{\text{LU}} \rho' \quad \Leftrightarrow \quad \rho' = (U_1 \otimes U_2 \otimes \cdots \otimes U_q) \rho (U_1 \otimes U_2 \otimes \cdots \otimes U_q)^{-1}$$

LU INVARIANTS FROM COLORED GRAPHS

From now on, restrict to pure state:

$$|\psi\rangle = \sum_{a_1, a_2, \dots, a_q} \underbrace{T_{a_1 a_2 \dots a_q}}_{\text{tensor}} |a_1\rangle \otimes |a_2\rangle \otimes \dots \otimes |a_q\rangle$$

Represent tensors and tensor contractions by *q-colored graphs*:

$$T_{a_1 a_2 \dots a_q} = \begin{array}{c} \bullet \\ \swarrow \text{red} \quad \downarrow \text{blue} \quad \searrow \text{green} \\ a_1 \quad a_2 \quad \dots \quad a_q \end{array} \quad \sum_c T_{abc} \overline{T_{cde}} = \begin{array}{c} \bullet \text{ (green)} \\ \swarrow \text{red} \quad \downarrow \text{blue} \quad \searrow \text{red} \\ a \quad b \quad d \quad e \end{array}$$

Any *closed q-colored graph* \mathcal{B} represents a *LU-invariant polynomial*, which we denote $\text{Tr}_{\mathcal{B}}(\overline{T}, T)$.

$$\begin{array}{c} \text{1} \\ \swarrow \text{blue} \quad \searrow \text{red} \\ \bullet \\ \swarrow \text{red} \quad \searrow \text{blue} \\ \text{2} \end{array} \leftrightarrow \overline{T_{abc}} T_{abc} \quad \begin{array}{c} \text{3} \\ \swarrow \text{green} \quad \searrow \text{green} \\ \bullet \\ \swarrow \text{red} \quad \searrow \text{red} \\ \text{2} \end{array} \leftrightarrow \overline{T_{abc}} T_{adc} \overline{T_{edf}} T_{ebf}$$

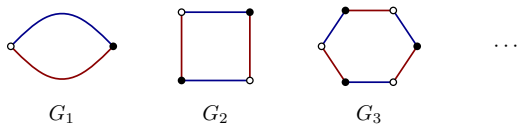
Claim. $\{\text{Tr}_{\mathcal{B}}(\overline{T}, T) \mid \mathcal{B} \text{ connected}\}$ generates the ring of polynomial LU-invariants. Any two $\mathcal{B}_1 \neq \mathcal{B}_2$ produce independent invariants in the limit of large dimension ($D_1, \dots, D_q \rightarrow +\infty$).

[R. Gurau's talk]

ENTANGLEMENT SPECTRUM OF A BIPARTITE SYSTEM

$$|\psi\rangle = \sum_{a_1, a_2} \underbrace{T_{a_1 a_2}}_{\text{rectangular matrix}} |a_1\rangle \otimes |a_2\rangle$$

\exists a single connected invariant of order $2n$, represented by cyclic graph G_n :



$$\text{Tr}_{G_n}(\bar{T}, T) = \text{tr}((TT^\dagger)^n) = \text{tr}(\rho_1^n) = \sum_i p_i^n$$

Entanglement structure of T characterized by its *entanglement spectrum*:

$$\begin{aligned} \{\text{LU invariants of } |\psi\rangle\} &\Leftrightarrow \{\text{tr}((TT^\dagger)^n)\}_{1 \leq n \leq \min(D_1, D_2)} \\ &\Leftrightarrow \text{Spec}(TT^\dagger) = \{p_i\}_{1 \leq i \leq \min(D_1, D_2)} \\ &\Leftrightarrow \{\text{singular values of } T\} = \{\sqrt{p_i}\}_{1 \leq i \leq \min(D_1, D_2)} \end{aligned}$$

Remark. Singular value decomposition \Rightarrow "Schmidt decomposition" :

$$|\psi\rangle = \sum_{i=1}^{\min(D_1, D_2)} \sqrt{p_i} |e_i\rangle \otimes |f_i\rangle$$

ENTANGLEMENT ENTROPIES

Functions of the entanglement spectrum which have special properties
(e.g. monotonicity under quantum operations) [P. Vrana's talk]

Entanglement entropy. *Von Neumann entropy* of ρ_1 (or ρ_2)

$$S(\rho_1) = -\text{tr}(\rho_1 \ln(\rho_1)) = - \sum_{i=1}^{\min(D_1, D_2)} p_i \ln(p_i)$$

Rényi- n quantum entropy. *Rényi- n entropy* ρ_1 (or ρ_2)

$$S_n(\rho_1) = \frac{1}{1-n} \ln(\text{tr}(\rho_1^n)) = \frac{1}{1-n} \ln \left(\sum_{i=1}^{\min(D_1, D_2)} p_i^n \right)$$

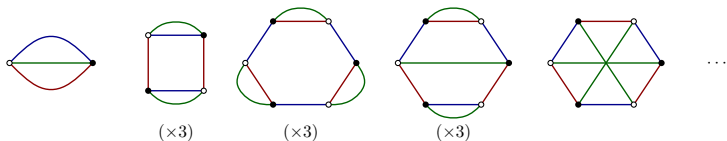
Example. Maximally entangled state / "Bell state"

$$|\psi\rangle = \frac{1}{\sqrt{D}} \sum_{i=1}^D |i\rangle \otimes |i\rangle \Rightarrow S_n(|\psi\rangle) = \ln D$$

→ flat entanglement spectrum.

MULTIPARTITE ENTANGLEMENT AND TENSOR INVARIANTS

$$|\psi\rangle = \sum_{a_1, a_2, \dots, a_q} \underbrace{T_{a_1 a_2 \dots a_q}}_{\text{tensor}} |a_1\rangle \otimes |a_2\rangle \otimes \dots \otimes |a_q\rangle, \quad \boxed{q \geq 3}$$



$\#\{\text{connected invariants of order } 2n\} = 1, 3, 7, 26, 97, 624, 4163 \dots$

Super-exponential growth of the number of independent invariants!

[Ben Geloun, Ramgoolam '13]

Question

Which of those many invariants are most relevant to characterize the multipartite entanglement structure of a many-body system?

TENSOR INVARIANTS AS PERMUTATIONS

Invariants can be conveniently parametrized by q -uplet of permutations

$$(\tau_1, \tau_2, \dots, \tau_q) \in S_n^{\times q}$$

[Ben Geloun, Ramgoolam '13]

$$\boxed{\{q\text{-colored graph of order } 2n\} \xleftrightarrow[1-1]{} S_n \backslash S_n^{\times q} / S_n}$$



$(n = 2, q = 3)$

$$\vec{\tau} = (\text{id}, \text{id}, \text{id})$$

$$\vec{\tau} = ((12), \text{id}, \text{id})$$

With this convention in mind, we will write:

$$\text{Tr}_{\mathcal{B}}(\bar{T}, T) = \text{Tr}_{\vec{\tau}}(\bar{T}, T)$$

Remark. $(q = 2) \text{tr}(\rho_1^n) = \text{Tr}_{G_n}(\bar{T}, T) = \text{Tr}_{(\text{id}, \sigma)}(\bar{T}, T)$, $\sigma := (12 \cdots n)$.

"Replica trick"

Review of some standard results on RTNs

1. Page curve
2. Area laws
3. RTNs with holographic area laws

Review of some standard results on RTNs

1. Page curve
2. Area laws
3. RTNs with holographic area laws

RANDOM BIPARTITE STATE: DEFINITION

$$|\psi\rangle = \sum_{i,j} M_{ij} |i\rangle \otimes |j\rangle \in \mathcal{H}_1 \otimes \mathcal{H}_2$$

Two equivalent ways of defining a *uniform random state*:

1. $|\psi\rangle = U|0\rangle$ with U random Haar-distributed on $U(D_1 D_2)$
2. M_{ij} Gaussian random matrix with covariance

$$\langle M_{ij} \bar{M}_{kl} \rangle = \frac{1}{D_1 D_2} \delta_{ik} \delta_{jl}$$

so that $\langle \text{tr}(MM^\dagger) \rangle = 1$.

For simplicity, let us adopt the second option, assuming $D_1 D_2 \gg 1$ (no further normalization required in this regime).

Goal: compute $\langle \text{tr}(\rho_1^n) \rangle = \langle \text{Tr}_{(\text{id}, \sigma)}(\bar{T}, T) \rangle$, and deduce typical entanglement spectrum / entanglement entropies.

RANDOM BIPARTITE STATE: $D_1 \ll D_2$ REGIME

Recall Wick's formula:

$$\langle M_{i_1 j_1} \bar{M}_{k_1 l_1} M_{i_2 j_2} \bar{M}_{k_2 l_2} \cdots M_{i_n j_n} \bar{M}_{k_n l_n} \rangle = \frac{1}{(D_1 D_2)^n} \sum_{\tau \in S_n} \prod_{p=1}^n \delta_{i_p k_{\tau(p)}} \delta_{j_p l_{\tau(p)}}$$

To simplify our life even more, let us first assume $D_1 \ll D_2$.

$$\begin{aligned} \langle \text{tr}(\rho_1^2) \rangle &= \langle \text{Tr}_{\square}(\bar{M}, M) \rangle = \langle \text{diag} \rangle = \langle \text{diag}_{\text{blue}} \rangle + \langle \text{diag}_{\text{red}} \rangle \\ &= \frac{1}{(D_1 D_2)^2} (D_1 D_2^2 + D_1^2 D_2) \approx \frac{1}{D_1} \end{aligned}$$

More generally, there is a single dominant Wick contraction for G_n :

$$\langle \text{tr}(\rho_1^n) \rangle \approx \frac{1}{(D_1 D_2)^n} D_1 D_2^n = D_1^{1-n}$$

RANDOM BIPARTITE STATE: $D_1 \ll D_2$ REGIME

Here, one can prove that

$$\ln \langle \text{tr}(\rho_1^n) \rangle \approx \langle \ln \text{tr}(\rho_1^n) \rangle$$

so that

$$\langle S_n(\rho_1) \rangle \approx \frac{1}{1-n} \ln(D_1^{1-n}) \approx \ln(D_1)$$

Taking the limit $n \rightarrow 1$ (can be justified...), we deduce

$$\boxed{\langle S(\rho_1) \rangle \approx \ln(D_1)}$$

Volume law

Is this behaviour more generally valid for $D_1 \leq D_2$?

RANDOM BIPARTITE STATE: $D_1 \leq D_2$ REGIME

$$\langle \text{tr}(\rho_1^n) \rangle = \langle \text{Tr}_{(\text{id}, \sigma)}(\bar{T}, T) \rangle \approx \frac{1}{(D_1 D_2)^n} \sum_{\tau \in S_n} D_1^{C(\tau)} D_2^{C(\tau^{-1} \circ \sigma)}$$

where

- ▶ $C(\tau) = \#$ cycles in the cycle-decomposition of τ ;
- ▶ $\sigma = (12 \cdots n)$.

In the regime $D_1 D_2 \rightarrow +\infty$ with $\frac{D_1}{D_2}$ fixed, the leading contributions maximize $C(\tau) + C(\tau^{-1} \circ \sigma)$.

Claim. This happens when

$$C(\tau) + C(\tau^{-1} \circ \sigma) = n + 1$$

Proof. $C(\tau) + C(\tau^{-1} \circ \sigma)$ counts the number of faces of a *combinatorial map* (discrete surface) with n edges and 1 vertex. We must therefore have

$$C(\tau) + C(\tau^{-1} \circ \sigma) = n + 1 - 2g$$

where $g \geq 0$ is the genus of the surfaces. □

RANDOM BIPARTITE STATE: $D_1 \leq D_2$ REGIME

Def. A *non-crossing permutation* is a $\tau \in S_n$ obeying:

$$C(\tau) + C(\tau^{-1} \circ \sigma) = n + 1$$

$$\Leftrightarrow \boxed{d(\text{id}, \tau) + d(\tau, \sigma) = d(\text{id}, \sigma)}$$

where $d(\tau_1, \tau_2) := n - C(\tau_1 \circ \tau_2^{-1})$ is the *Cayley distance* on S_n . One also says that τ is *on a geodesic* between id and σ .

With these definitions, one finds

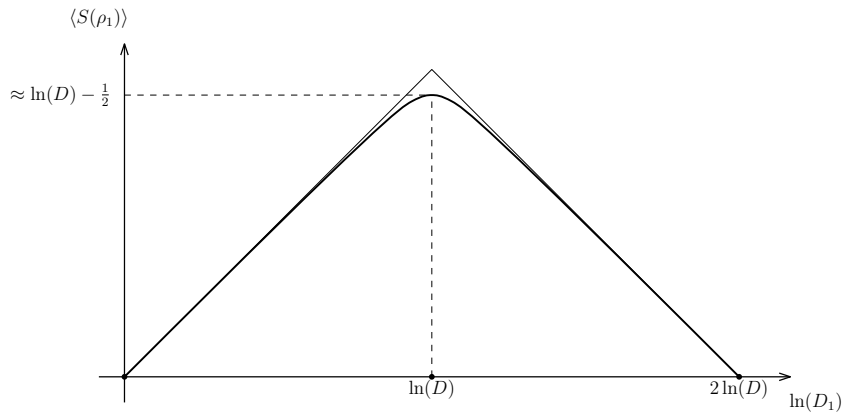
$$\langle \text{tr}(\rho_1^n) \rangle \approx D_1^{1-n} \sum_{\tau \in S_n \text{ n.c.}} \left(\frac{D_1}{D_2} \right)^{C(\tau)-1} = D_1^{1-n} {}_2F_1(1-n, -n, 2; \frac{D_1}{D_2})$$

After analytic continuation of n , we can compute $\lim_{n \rightarrow 1} \frac{1}{1-n} \ln \langle \text{tr}(\rho_1^n) \rangle$, yielding

$$\boxed{\langle S(\rho_1) \rangle \approx \underbrace{\ln(D_1)}_{\text{volume law}} - \underbrace{\frac{D_1}{2D_2}}_{\text{finite correc.}}}$$

RANDOM BIPARTITE STATE: PAGE CURVE

Set $D^2 := D_1 D_2$.



[Page '93; Foong, Kano '94; Sánchez-Ruiz '95; Sen '96]

Review of some standard results on RTNs

1. Page curve
2. Area laws
3. RTNs with holographic area laws

CONDENSED MATTER

[TALKS BY N. SCHUCH AND M. C. BAÑULS]

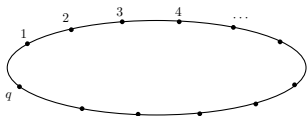
For a random $|\psi\rangle \in \mathcal{H}_A \otimes \mathcal{H}_{\bar{A}}$:

$$\langle S(\rho_A) \rangle \sim \ln(d_{\text{loc}}^{V_A}) \sim V_A$$

However, the ground state of a **gapped local Hamiltonian** typically obeys an **area law**

$$S(\rho_A) \leq K|\partial A|$$

→ **Tensor Networks (TN)**: variational Ansatz with a **polynomial number of parameters** in system size, which obey area laws by design.

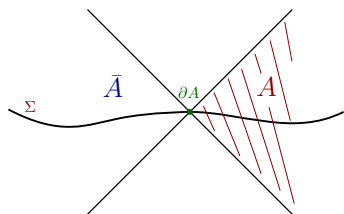


Matrix Product State

$$|\psi\rangle = \sum_{a_1, a_2, \dots, a_q} \text{tr}(M^{a_1} M^{a_2} \dots M^{a_q}) |a_1\rangle \otimes \dots \otimes |a_q\rangle$$

RELATIVISTIC QFT

Unruh effect / Bisognano-Wichman theorem (and, similarly, BH entropy)



$$|\rho_A\rangle = \frac{e^{-2\pi H_A}}{Z}$$

with H_A = boost Hamiltonian.

Universal divergence

$$S(\rho_A) \propto K|\partial A|$$

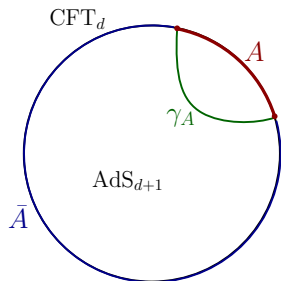
Suggested (heuristic) entropic derivation of general relativity [Jacobson '95, '15]

$$\left\{ \begin{array}{l} \delta S(\rho_A) = \underbrace{\eta}_{\text{universal}} \delta|\partial A| \\ \text{"entanglement equilibrium hypothesis"} \end{array} \right.$$

$$\iff G_{ab} + \Lambda g_{ab} = 8\pi G T_{ab}$$

HOLOGRAPHY

Holographic area law (AdS/CFT):



Ryu-Takayanagi formula

$$S(\rho_A) = \frac{|\partial A|}{4G}$$

Holographic version of Jacobson's argument:

Ryu-Takayanagi for any ball $A \Rightarrow$ Einstein's equations up to 2nd order

[Faulkner et al. 2017]

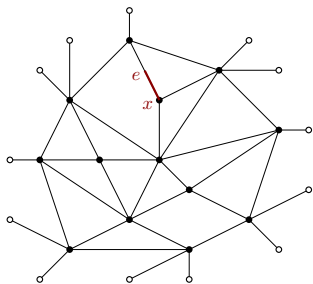
Review of some standard results on RTNs

1. Page curve
2. Area laws
3. RTNs with holographic area laws

DEFINITION

[M. WALTER'S TALK]

Graph $G = (V, E)$ with $V = V_{\text{bulk}} \sqcup V_{\partial} = \{\bullet\} \sqcup \{\circ\}$



Associate Hilbert space $\mathcal{H}_{e,x}$ to each half-edge (e, x) , with

$$\dim \mathcal{H}_{e,x} = D \quad (\text{bond dimension})$$

Edge data. Maximally mixed bipartite state along each edge

$$|\phi\rangle = \bigotimes_{(x,y) \in E} \frac{1}{\sqrt{D}} \sum_{i=1}^D |i\rangle_x \otimes |i\rangle_y$$

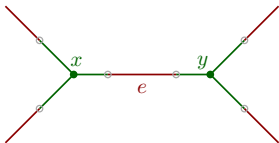
Vertex data.

$$|\eta\rangle = \bigotimes_{x \in V_{\text{bulk}}} |\eta\rangle_x, \quad |\eta\rangle_x \in \bigotimes_{e \text{ incident to } x} \mathcal{H}_{e,x}$$

DEFINITION

Tensor Network state:

$$|\psi\rangle = (\text{Id}_{V_\partial} \otimes \langle \eta |) |\phi\rangle \in \bigotimes_{v \in V_\partial} \mathcal{H}_{e_v, v}$$

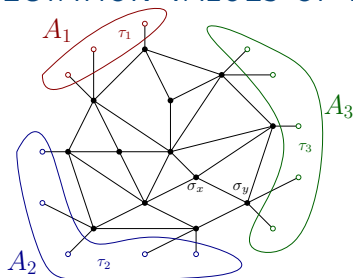


Random Tensor Network state. Take $\{|\eta\rangle_x\}$ to be independent, Haar-distributed (or Gaussian) random vectors. We then have the Wick formula

$$\mathbb{E} (|\eta\rangle_x \langle \eta|^{\otimes n}) \propto \sum_{\sigma \in S_n} R_x(\sigma)$$

where $R_x(\sigma)$ is a permutation operator acting on $\mathcal{H}_x^{\otimes n}$

EXPECTATION VALUES OF LU INVARIANTS



Define

$$\mathcal{Z}(\tau_1, \tau_2, \dots, \tau_q) := \langle \text{Tr}(\tau_1, \tau_2, \dots, \tau_q)(\bar{T}, T) \rangle$$

Partition function of generalized spin model

$$\mathcal{Z}(\tau_1, \tau_2, \dots, \tau_q) = \sum_{\substack{\sigma: V \rightarrow S_n \\ \forall v \in A_k, \sigma(v) = \tau_k}} \exp(-\beta \mathcal{E}(\sigma))$$

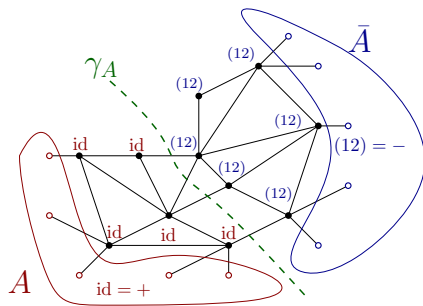
$$\mathcal{E}(\sigma) = \sum_{(x,y) \in E} d(\sigma(x), \sigma(y)) \quad \beta = \ln(D)$$

ENTROPY \sim GROUND STATE ENERGY

The $D \rightarrow +\infty$ limit is a *zero-temperature limit*. At leading-order, we thus have:

$$\mathcal{Z}_{(\tau_1, \tau_2, \dots, \tau_q)} = \mathcal{N}_{g.s.} D^{-\mathcal{E}_{\min}} (1 + \mathcal{O}(1/D))$$

Example. $\mathcal{Z}_{(\text{id}, (12))} = \langle \text{tr}(\rho_A^2) \rangle$ is a Ising partition function.



Let γ_A be a *min-cut*, that is: a minimal collection of edges that disconnect A from \bar{A} when cut.

$$\mathcal{E}_{\min} = |\gamma_A|$$

$$\langle \text{tr}(\rho_A^2) \rangle \approx \mathcal{N}_{g.s.} D^{-|\gamma_A|}$$

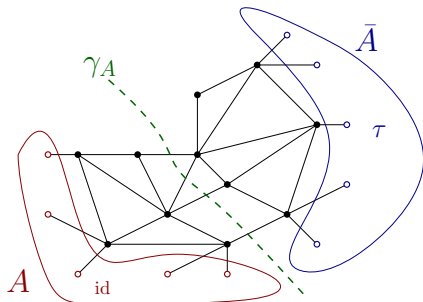
Holographic area law for Rényi-2 entropy:

$$\langle S_2(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

BIPARTITE SETTING: GENERAL RESULTS

Min-cut / Max-flow theorem. (simplified version)

There exists $|\gamma_A|$ edge-disjoint paths connecting A to \bar{A} .



Proof. Explicit algorithm: Ford-Fulkerson (50's).

Theorem. We have $\mathcal{Z}_{(id,\tau)} \approx \mathcal{N}_{g.s.} D^{-d(id,\tau)|\gamma_A|}$. In particular, taking $\tau = (12 \cdots n)$ leads to

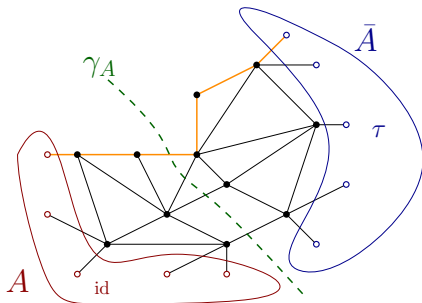
$$\langle S_n(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

[Collins, Nechita, Zyczkowski '13 ; Hayden et al '16; ...]

BIPARTITE SETTING: GENERAL RESULTS

Min-cut / Max-flow theorem. (simplified version)

There exists $|\gamma_A|$ edge-disjoint paths connecting A to \bar{A} .



Proof. Explicit algorithm: Ford-Fulkerson (50's).

Theorem. We have $\mathcal{Z}_{(id, \tau)} \approx \mathcal{N}_{g.s.} D^{-d(id, \tau) |\gamma_A|}$. In particular, taking $\tau = (12 \cdots n)$ leads to

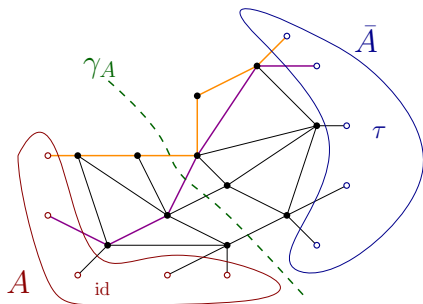
$$\langle S_n(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

[Collins, Nechita, Zyczkowski '13 ; Hayden et al '16; ...]

BIPARTITE SETTING: GENERAL RESULTS

Min-cut / Max-flow theorem. (simplified version)

There exists $|\gamma_A|$ edge-disjoint paths connecting A to \bar{A} .



Proof. Explicit algorithm: Ford-Fulkerson (50's).

Theorem. We have $\mathcal{Z}_{(id,\tau)} \approx \mathcal{N}_{g.s.} D^{-d(id,\tau)|\gamma_A|}$. In particular, taking $\tau = (12 \cdots n)$ leads to

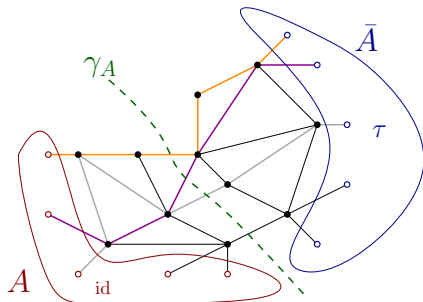
$$\langle S_n(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

[Collins, Nechita, Zyczkowski '13 ; Hayden et al '16; ...]

BIPARTITE SETTING: GENERAL RESULTS

Min-cut / Max-flow theorem. (simplified version)

There exists $|\gamma_A|$ edge-disjoint paths connecting A to \bar{A} .



Proof. Explicit algorithm: Ford-Fulkerson (50's).

Theorem. We have $\mathcal{Z}_{(id,\tau)} \approx \mathcal{N}_{g.s.} D^{-d(id,\tau)|\gamma_A|}$. In particular, taking $\tau = (12 \cdots n)$ leads to

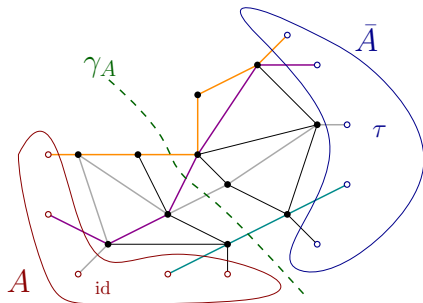
$$\langle S_n(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

[Collins, Nechita, Zyczkowski '13 ; Hayden et al '16; ...]

BIPARTITE SETTING: GENERAL RESULTS

Min-cut / Max-flow theorem. (simplified version)

There exists $|\gamma_A|$ edge-disjoint paths connecting A to \bar{A} .



Proof. Explicit algorithm: Ford-Fulkerson (50's).

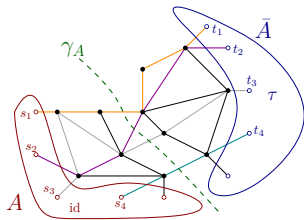
Theorem. We have $\mathcal{Z}_{(id,\tau)} \approx \mathcal{N}_{g.s.} D^{-d(id,\tau)|\gamma_A|}$. In particular, taking $\tau = (12 \cdots n)$ leads to

$$\langle S_n(\rho_A) \rangle \approx |\gamma_A| \ln(D)$$

[Collins, Nechita, Zyczkowski '13 ; Hayden et al '16; ...]

BIPARTITE SETTING: GENERAL RESULTS

Proof. Repeated application of the triangular inequality along each flow.

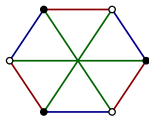
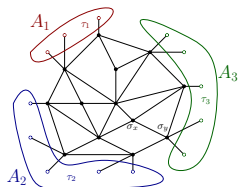


Let $\mathcal{P} = \mathcal{P}_1 \sqcup \dots \sqcup \mathcal{P}_{|\gamma_A|}$ the edges of a max-flow.

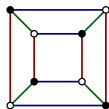
$$\begin{aligned} \mathcal{E}(\sigma) &= \sum_{(x,y) \in E} d(\sigma(x), \sigma(y)) \geq \sum_{i=1}^{|\gamma_A|} \sum_{(x,y) \in \mathcal{P}_i} d(\sigma(x), \sigma(y)) \\ &\geq \sum_{i=1}^{|\gamma_A|} d(\sigma(s_i), \sigma(t_i)) = |\gamma_A| d(\text{id}, \tau) \end{aligned}$$

And those inequalities are saturated. □

MULTIPARTITE SETTING: TWO EXAMPLES



$(\text{id}, (123), (132))$



$((12)(34), (13)(24), (14)(23))$

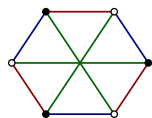
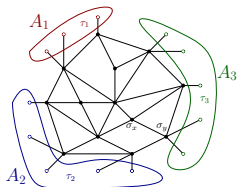
Definition. (τ_1, τ_2, τ_3) is *compatible* if $\exists \sigma$ such that:

- ▶ $d(\tau_1, \sigma) + d(\sigma, \tau_2) = d(\tau_1, \tau_2)$;
- ▶ $d(\tau_1, \sigma) + d(\sigma, \tau_3) = d(\tau_1, \tau_3)$;
- ▶ $d(\tau_2, \sigma) + d(\sigma, \tau_3) = d(\tau_2, \tau_3)$.

Claim.

- ▶ $(\text{id}, (123), (132))$ is *compatible*;
- ▶ $((12)(34), (13)(24), (14)(23))$ is *not compatible*.

MULTIPARTITE SETTING: TWO EXAMPLES



$(\text{id}, (123), (132))$

Theorem. For $\vec{\tau} = (\text{id}, (123), (132))$

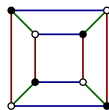
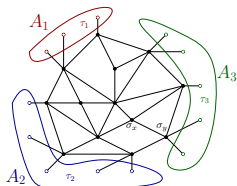
$$\mathcal{E}_{\min} = |\gamma_{A_1}| + |\gamma_{A_2}| + |\gamma_{A_3}|$$

Proof. Generalization of previous argument relying on *multi-cut / max-flow theorem*... □

[Cui, Hayden, et al '18 ; Dong, Qi, Walter '21; Kudler-Flam, Ryu, Narovlansky '21]

→ Entanglement measure which, at leading order, does not capture information that was not already contained in bipartite measures.

MULTIPARTITE SETTING: TWO EXAMPLES



$((12)(34), (13)(24), (14)(23))$

Theorem. For $\vec{\tau} = ((12)(34), (13)(24), (14)(23))$

$$\mathcal{E}_{\min} = 4 \min_{\text{tripartition } \gamma} |\gamma|$$

[Penington, Walter, Witteveen '22]

→ Entanglement measure capturing genuinely tripartite information.

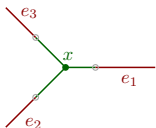
Can we explore the space of multipartite entanglement measures more systematically? at least for simple networks?

[w.i.p with Johann Chevrier, Luca Lionni, Michael Walter]

RTNs with reduced randomness: local
Haar-averaging

BASIC IDEA

Keep local entanglement structure at each vertex of the network fixed i.e. *average over local unitaries* instead of full unitary group.



$$\mathcal{H}_x = \mathcal{H}_{x,e_1} \otimes \mathcal{H}_{x,e_2} \otimes \mathcal{H}_{x,e_3}$$

Questions:

- ▶ Can such reduced randomness support area laws?
- ▶ If so, can we identify specific entanglement structures which do so?
- ▶ Can this framework produce richer entanglement spectra than in the standard case (non-flat spectra)?

LOCAL HAAR-AVERAGING

Suppose x is a q -valent vertex, $q \geq 3$.

$$|T\rangle_x = \sum_{a_1, a_2, \dots, a_q} \underbrace{T_{a_1 a_2 \dots a_q}}_{\text{fixed tensor}} |a_1\rangle \otimes |a_2\rangle \otimes \dots \otimes |a_q\rangle$$

Average over $U(D)^{\otimes q} \not\subseteq U(Dq)$:

$|\Psi\rangle_x \sim$ uniform random state in LU-orbit of $|T\rangle_x$, that is:

$$\mathbb{E} [|\Psi\rangle_{xx} \langle \Psi|^{\otimes n}] := \int_{U(D)^{\otimes q}} dU \left[\left(\bigotimes_{c=1}^q U^{(c)} \right) \underbrace{|T\rangle_{xx} \langle T|}_{\text{seed state}} \left(\bigotimes_{c=1}^q U^{(c)} \right)^\dagger \right]^{\otimes n}$$

EXPLICIT EVALUATION OF MOMENTS

Weingarten calculus allows to write:

[Collins et al. '00s]

$$\mathbb{E} [|\Psi\rangle_{XX}\langle\Psi|^{\otimes n}]_{\{i_s, j_s\}_{s=1}^n} = \sum_{\sigma} \underbrace{F_T(\sigma)}_{\text{state-dep. weight}} \underbrace{\mathcal{I}_{\{i_s, j_s\}_{s=1}^n}^{\sigma}}_{\substack{\text{tensor structure} \\ \sim \text{colored diagram}}}$$

where $\sigma = (\sigma_1, \dots, \sigma_q) \in S_n^{\times q}$ etc. and

$$\mathcal{I}_{\{i_s, j_s\}_{s=1}^n}^{\sigma} = \prod_{c=1}^q \prod_{s=1}^n \delta_{i_s^c, j_{\sigma_c(s)}^c}$$
$$F_T(\sigma) = \sum_{\tau} \underbrace{\text{Tr}_{\tau}(\bar{T}, T)}_{\text{LU invariant}} \underbrace{\mathbf{W}^D(\sigma\tau^{-1})}_{\substack{\text{Weingarten} \\ \text{functions}}}$$

Instead of one permutation per vertex, we have **two multiplets σ and τ** .

However, the Weingarten functions encourages σ and τ to be "close to each other".

ASYMPTOTICS OF THE WEINGARTEN FUNCTION

$$D^n W^{(D)}(\sigma\tau^{-1}) = D^{-\overbrace{d(\sigma, \tau)}^{\text{Cayley distance}}} \underbrace{M(\sigma\tau^{-1})}_{\substack{\text{Moebius function} \\ \text{(no definite sign)}}} (1 + \mathcal{O}(1/D^2))$$

For some entangled seed states, this allows to write:

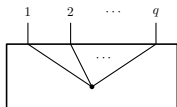
$$F_T(\boldsymbol{\sigma}) = \sum_{\boldsymbol{\tau}_1} D^{-\kappa(\mathcal{B}_{\boldsymbol{\tau}}) - \sum_{c=1}^q d(\sigma_c, \tau_c) - \omega_T(\boldsymbol{\tau})} \left(\prod_{c=1}^q M(\sigma_c \tau_c^{-1}) \right) (1 + \mathcal{O}(1/D^2))$$

where

- ▶ $\kappa(\mathcal{B}_{\boldsymbol{\tau}}) = n - \#\{\text{connected components of } \mathcal{B}_{\boldsymbol{\tau}}\}$ is minimal when $\tau_1 = \tau_2 = \dots = \tau_q$;
- ▶ $\sum_{c=1}^q d(\sigma_c, \tau_c)$ is minimal when $\boldsymbol{\sigma} = \boldsymbol{\tau}$;
- ▶ $\omega_T(\boldsymbol{\tau})$ is a state-dependent contribution.

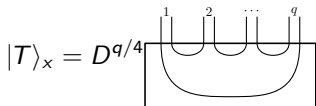
$\omega_T(\boldsymbol{\tau})$ FOR SOME EXAMPLES OF SEED STATES

1. GHZ state:



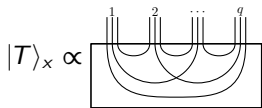
$$|T\rangle_x = D^{(q-1)/2} \sum_{i=1}^D \bigotimes_{c=1}^q |i\rangle_c \Rightarrow \boxed{\omega_T(\boldsymbol{\tau}) = 0}$$

2. "Cyclic" state:



$$|T\rangle_x = D^{q/4} \Rightarrow \boxed{\omega_T(\boldsymbol{\tau}) = \underbrace{g_J(\mathcal{B}_{\boldsymbol{\tau}})}_{\text{genus of "jacket" } J=(12\cdots q)}}$$

3. "Complete graph" state:



$$\Rightarrow \boxed{\omega_T(\boldsymbol{\tau}) = \omega_{\text{Gurau}}(\mathcal{B}_{\boldsymbol{\tau}})}$$

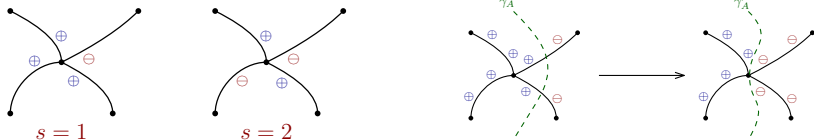
[R. Gurau's talk]

EXAMPLE 1: GHZ WITH $n = 2$ AND $q = 4$

$S_2 = \{\oplus, \ominus\}$. One can explicitly sum over τ , and derive a generalized spin model governed by energy:

$$\mathcal{E}(\{\sigma_v\}) = \mathcal{E}_{\text{Ising}}(\{\sigma_v\}) + \underbrace{2\nu_1 + \nu_2}_{\text{vertex defects}}$$

where $\nu_s := \#\{\text{defects of type } s\}$.



Claim. Energy minimizers are not Ising configurations in general; as a result: we have $c_A \leq |\gamma_A|$ (with $c_A < |\gamma_A|$ for some networks) such that

$$\mathbb{E}[\text{tr}(\rho_A^2)] = \mathcal{N}_{\text{g.s.}} \exp(-\ln(D)c_A) (1 + \mathcal{O}(1/D)) ,$$

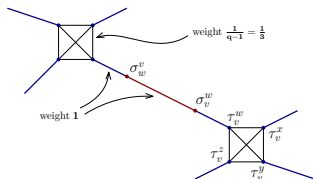
$$\mathbb{E}[S_2(\rho_A)] \approx c_A \ln(D)$$

EXAMPLE II: COMPLETE GRAPH SEED STATE

A "spin" configuration is labelled by

$$s = \{\sigma_v^w, \sigma_w^v, \tau_v^w, \tau_w^v \mid (v, w) \in E\}$$

Leading order contributions minimize the Ising energy $\mathcal{E}_{\text{Ising}}(s)$ of a *refined network*



Conjecture. The Rényi entropy of a subregion A is governed by the size of a minimal-cut $|\gamma_A|$

$$\mathbb{E}[\text{tr}(\rho_A^2)] = \mathcal{N}_{\text{g.s.}} \exp(-\ln(D)|\gamma_A|) (1 + \mathcal{O}(1/D)), \quad \boxed{\mathbb{E}[S_2(\rho_A)] \approx |\gamma_A| \ln(D)}$$

To be checked: absence of cancellation in the leading-order sector (which could arise due to non-positivity of Moebius function). Treacherous...

CONCLUSION

- ▶ RTNs with reduced randomness, allowing for tunable entanglement structure at each vertex.
- ▶ Rényi entropy evaluation maps to generalized spin models: permutation associated to half-edges rather than edges, and energy contribution from internal structure of vertices.
- ▶ First examples with homogeneous choice across the network suggest that distinct choices of local entanglement structures affect the entanglement spectrum of the global state.
- ▶ In principle, the local entanglement structure could be chosen non-homogeneously across the network → large variety of effective behaviour can be expected, but likely hard to investigate in detail.
- ▶ How is the multipartite entanglement structure of the global state affected?